CBCS SCHEME

USN 15CS65

Sixth Semester B.E. Degree Examination, Dec.2019/Jan.2020 Data Mining and Data Warehousing

Time: 3 hrs.

Note: Answer any FIVE full questions, choosing ONE full question from each module.

Module-1

a. What is Data warehouse? Explain three tier architecture of data warehouse.
b. Explain the schemas of multidimensional data models.
(08 Marks)
(08 Marks)

OR

a. What is Data cube measure? Explain the categorization of measures.
 b. Explain data cube operations with examples.
 (08 Marks)
 (08 Marks)

Module-2

a. Explain data cube computation and curse of dimensionality.
 b. Explain different methods of indexing OLAP data.

(08 Marks)
(08 Marks)

OR

- 4 a. State and explain various data mining tasks. (08 Marks)
 - b. Define Similarity and dissimilarity between the objects. Find SMC and Jaccord's coefficient of two binary vectors.
 X = (1, 0, 0, 0, 0, 0, 0, 0, 0)
 Y = (0, 0, 0, 0, 0, 1, 0, 0, 1).
 (08 Marks)

Module-3

What is Association Analysis? Explain Association rule, Support and Confidence. (08 Marks)
 State Apriori principle. Write apriori algorithm for frequent itemset. (08 Marks)

OR

6 a. Construct an FP tree for the following dataset.

TID	Items
1	{a, b}
2	{b, c, d}
3	{a, c, d, e}
4	{a, d, e}
5	{a, b, c}
6	{a, b, c, d}
7	{a}
8	{a, b, c}
9	{a, b, d}
10	{b, c, e}

b. Explain the strategies used in frequent itemset generation. (08 Marks)

(08 Marks)

Max. Marks: 80

Module-4

a. Explain the general approach for solving classification problem. (08 Marks)

Write the algorithm for decision tree induction. (08 Marks)

of 2

Important Note: 1. On completing your answers, compulsority draw diagonal cross lines on the remaining blank pages.

2. Any revealing of identification, appeal to evaluator and /or equations written eg, 42+8 = 50, will be treated as malpractice.

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15CS651

OR

a. Explain the methods of comparing classifiers.b. Write the characteristics of nearest neighbor classifier.

(08 Marks) (08 Marks)

Module-5

Explain the requirements of cluster analysis.

State and explain K - means algorithm.

(08 Marks (08 Marks

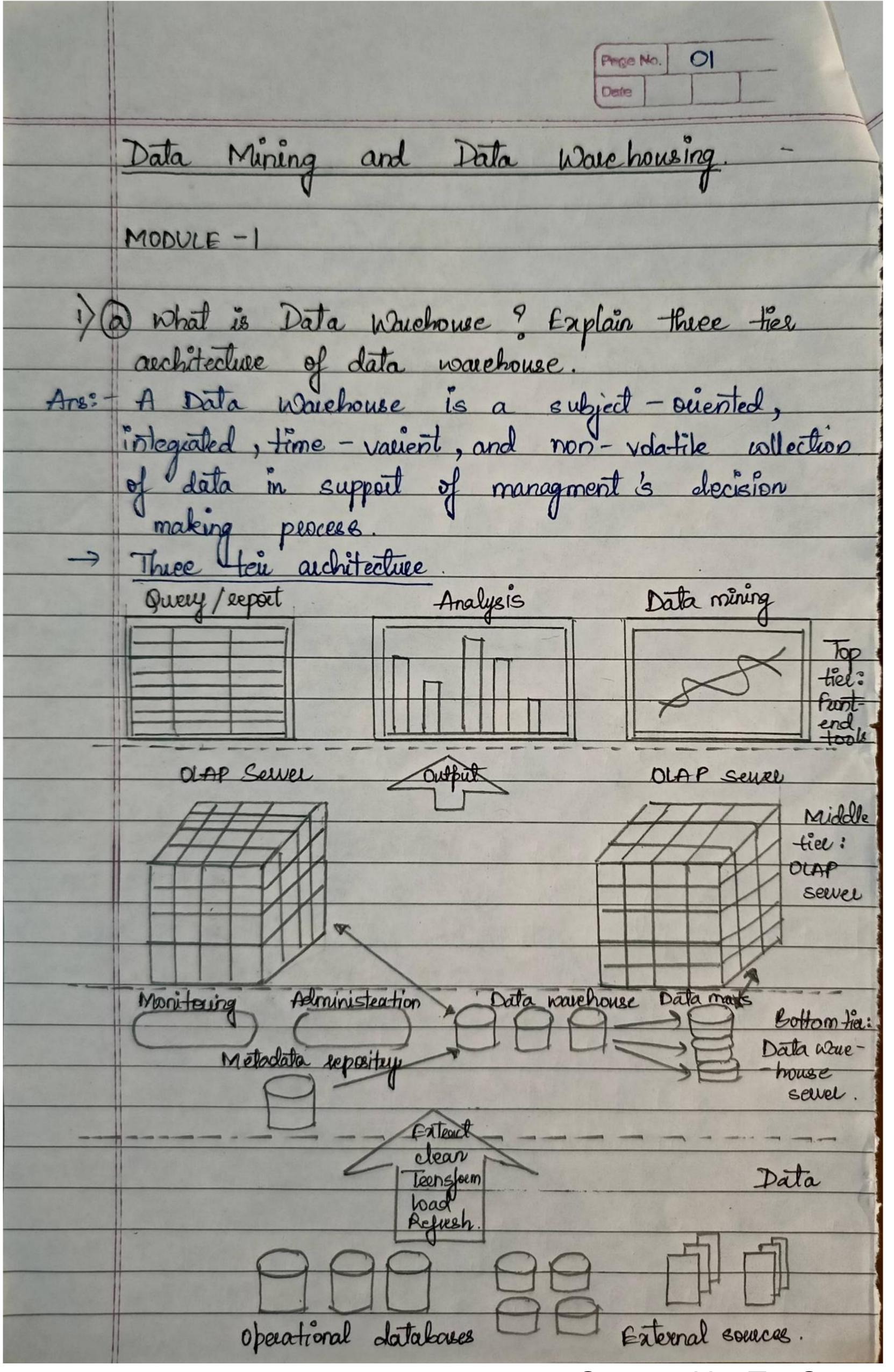
OR

10 a. Write DBSCAN clustering algorithm and estimate time and space complexity.

b. State and explain the issues in cluster evaluation.

(08 Marks) (08 Marks)

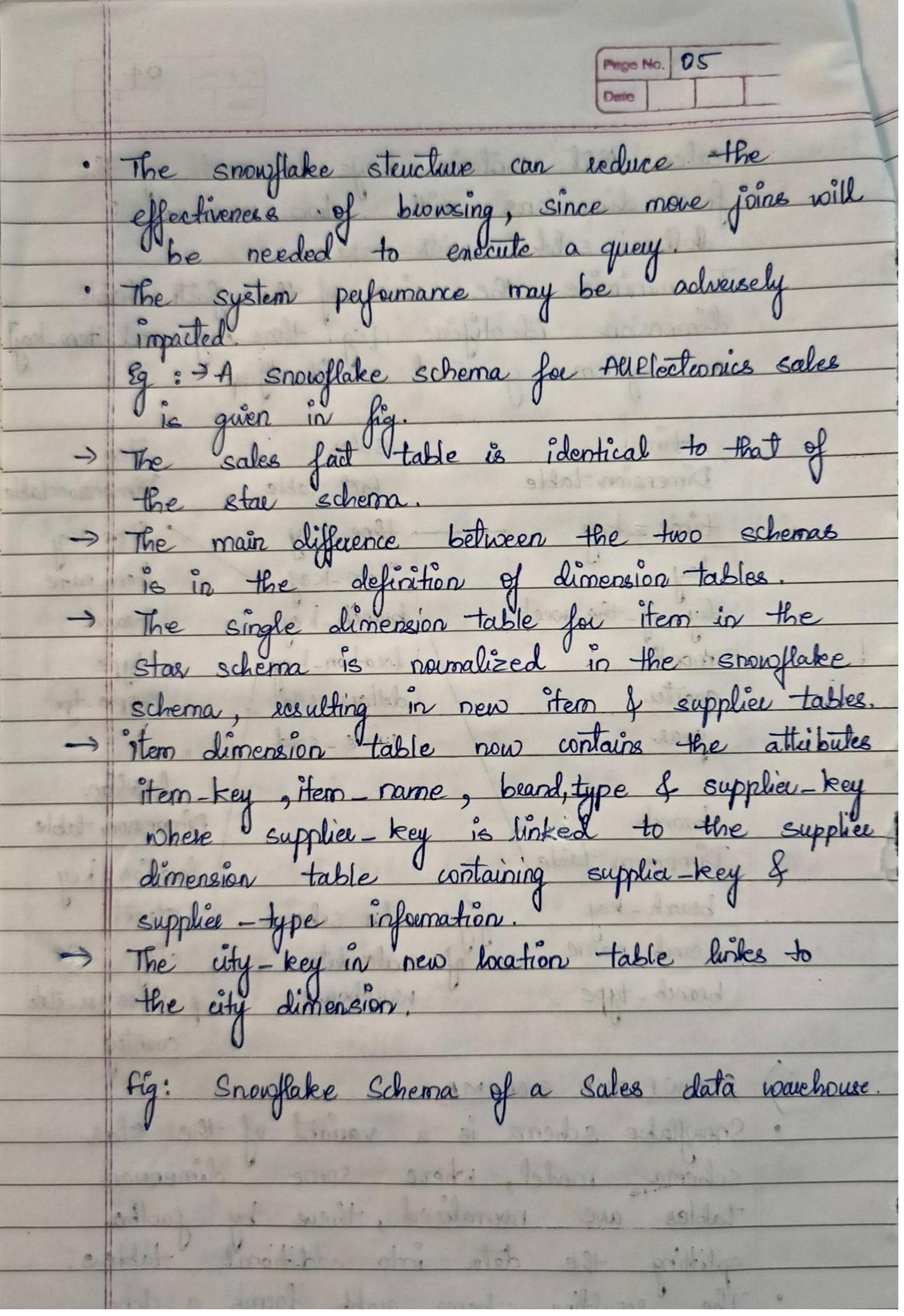
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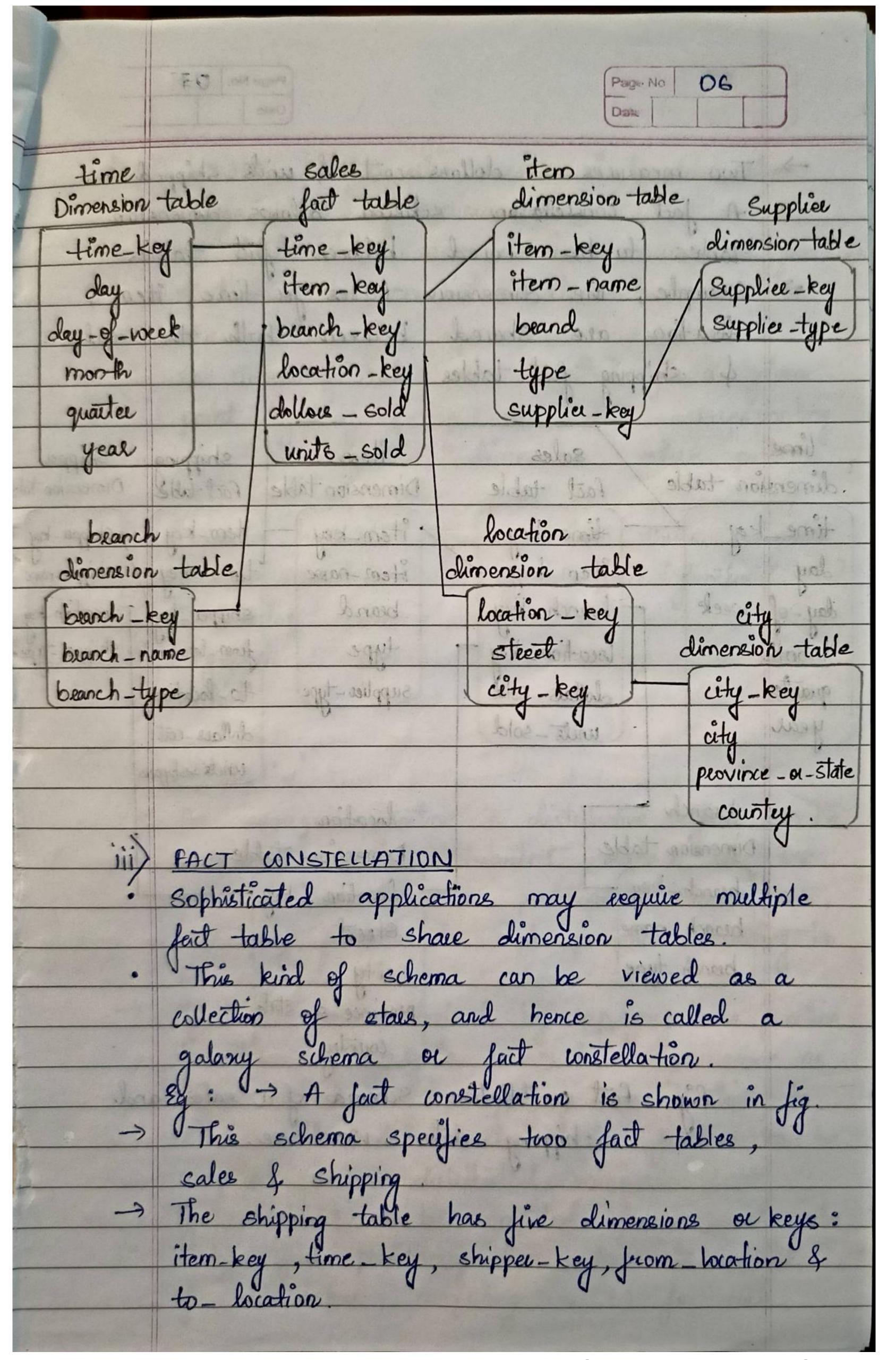


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	Prege No. 03 Dete
b)	Multidimensional DLAP (MOLAP) model i.e., a special purpose saves that discolly implements multidimensional
3203	data and operations.
. 6	The top tiee: - is a front-end client layer, which contains query and reporting tooks,
James 1	eg: teend analysis, prediction and so on.
1)6	Explain the schemas of multidimensional data models.
Ans:	Schemas for Multidimesional data models.
i	STAR SCHEMA
Soul of	The most commonly modelling paradigm is the Star Schema, in which the data warehouse
9	a large central table [fact table] containing the bulk of data, with no ecdundarcy
b).	a set of smaller attendant tables [dimension table],
	The schema graph ecsembles à starbuent, with the dimension tables displayed in a ladial
1.2+	pattern around the central fact table. Eg: - A stook schema for All Electronics sales
and:	Sales are considered along four dimensions: item,
	time, bearch of location.
	The schema contains a conteal fact table for

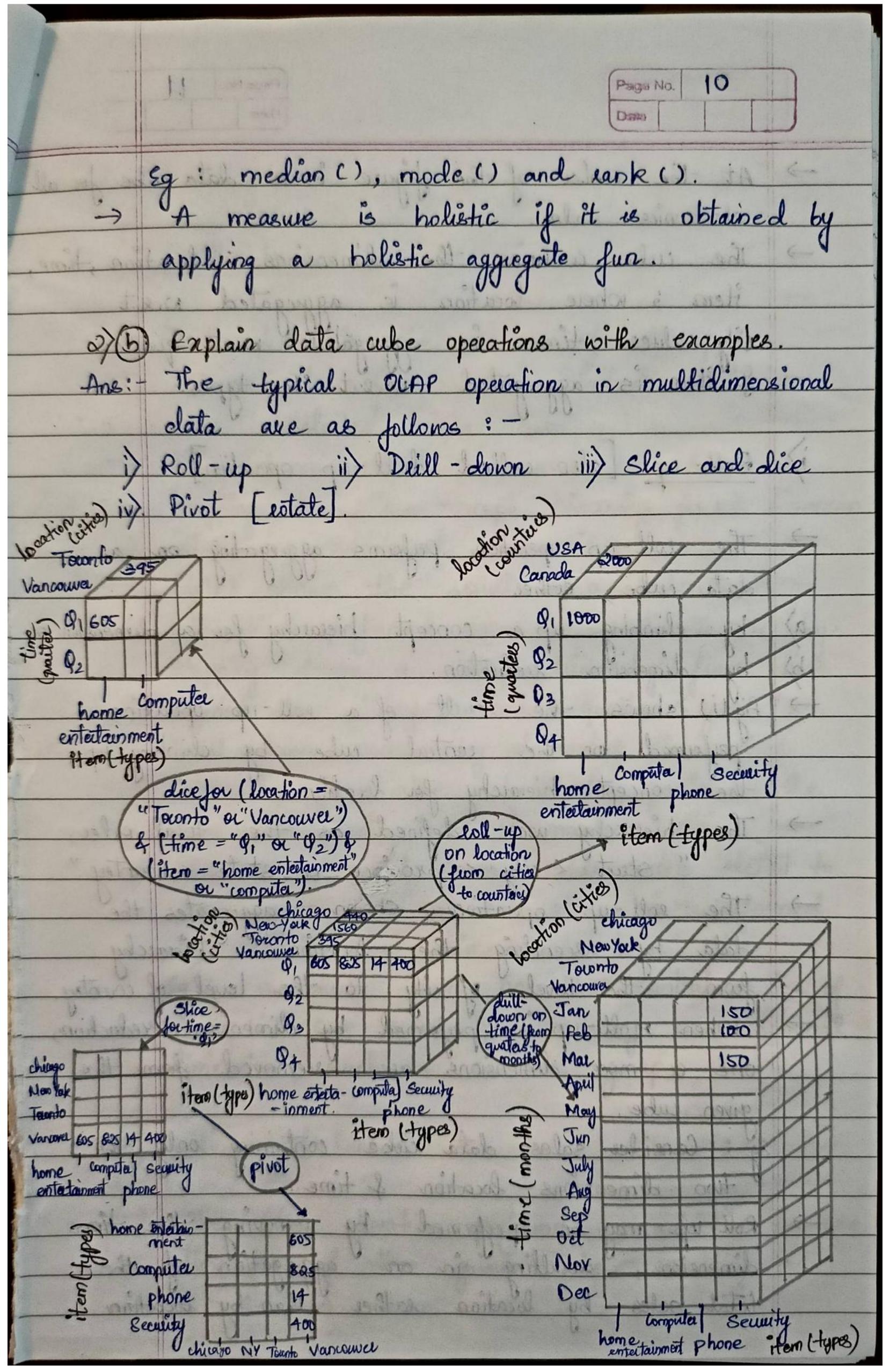
		Proges	No. 04
	sales that contain four dimensions	is keys to ear	ch of the
	dollars - sold & u	nits-sold.	
pi q	To minimize the	size of the	fait table
6	dimension identif	ies [eg. time-	egy and item-key
	are system genera	identifier	
	0 0	The same state of	
Je to	time	Sales	item
16	Dimension table	fact table	Dimension-table
30303	time_key	time_key	item-key
	day	item-key	item_name
	day - of - the-week	beanch-key	beand
	month	/ location-key	type
A MARKET N	quarter	dollars_sold	Supplia-type
adiation of	year /	units_sold	ab cost 16-
1000	wa de aquillond	good costi us	location
to day and	bearch	1000 - 100 -	Dimension table
1	Dimension table	ins tables to	location - key
	bearch - key	g: Star schema	steel
eb-	bearch - name	of sales data	city
	beanch-type	vouehouse.	per vince - or-state
			country
(ii)	SNOWFLAKE SCHEM	1A made valor	2 : 177
•	Snowflake schema i	s a variant of	the star
	schema model,	where some	dimension
	tables are noun	alized, there by	furthee
	splitting the data	into additiona	l tables.
•	The lesulting scho	ema geaph forms	
	similar to snow	lake.	
	V		





			Peop	e No. 07
			1 0	
→	Two measu	ues : dollars	_cost & une	ts_shipped
130	A fait	constellation s	chema allow	s dimension
-3 13.01 ENCS	tables to	be shared	between f	act tables.
1001-)			tables joi	time, item & oth the sales
111111111111111111111111111111111111111	location			orn the saves
	4 shippi	ng fact tables.	blos _ wallab	Jahrenp
time		Sales	item item	shipping shipper
dimens	ion table	fact table	Dimension table	
time_	key /	time legy	item_key	item-key shipper-key
day	5	Hem-key	Hem-name	time key / shipper-name
day-of	- week -	beanch-key	beand	shipper-key Thocation-key
month	pians mills	location_key 7	type	from-bocation. Shipper-type
quaitel	الما منابور	chothas - sold	Supplier-type	to-location
year	11-10	units-sold		dollars-cost
shitz 10 - 1	iscurosq -			unite-shippers
133	beanch -		location	
1	Dimension tab	le	Dimension to	ible as
state	beanch-key	x 1000 20015	- location - ke	y .
	beanch_name	asica-dimension	Steed	11+ 15-1-
	beanch-type	nd and ha	city	d sarv.
+	is called	and beace	province ou - sta	te atsolves
	Marion .	strine trak	country	Jundop
21	fig : fo	ict constellation	Schema of	a'sales and
6.	estilet to	chipping dat	ta vacchous	e. solle
		. 0	aringide J	l estes
2 1193	10 200 2000	than the oliv	significant solde	10 011 6
2 !!!	lecon brook	shipper key	und and	No d. moti
		"	2	

for each subsube. Hence, count () is distributive aggrégate function sur (), min () & man () are distributive aggregative functions. A measure is distributive if it is obtained by applying a distributive aggregate function. Algebraic An aggregate furction is algebeic, if it can be computed by an algebraic function with M arguments [where M is bounded positive integer] Leach of which is obtained by applying a distributive aggregate function. Eg: aug () [average] can be computed by sum()
where both sum() & count() are distributive aggrégate junctions. min_NC) and man_NC) and standard - deviation () are algebraic aggrégate functions. -) A measure is algebic if it is obtained by applying an algebeic eggregate functions. An aggregate function is holistic if there is no constant bound on the storage size needed to describe a subaggregate lie, there does not exist an algebric function with M arguments [where M is constant] that characterizes the Datistana)

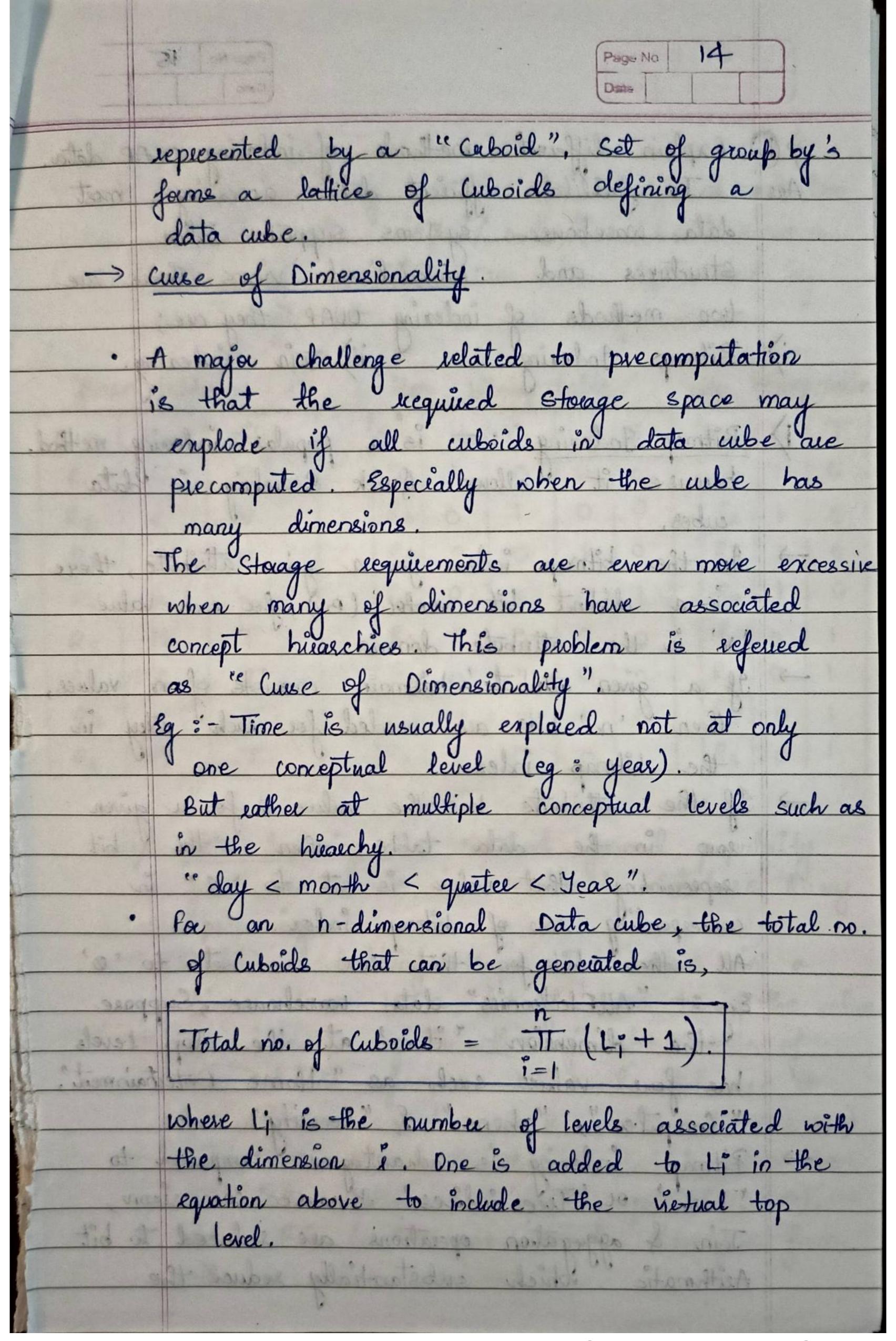


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->	At the center of the figure is a data cube for all
v3 1	At the center of the figure is a data cube for all electronics sales.
->	The cube contains the dimensions: location, time,
	Hern; where location is aggregated w.e.t
	sitem is aggregated w. e.t item types.
Jeneti	f item is aggregated w. e.t item types.
	a would do sue time
i)	Roll-up [also called dill-up operation]
	office I state I took to
-	The eall-up operation performs aggregating on a data cube, either
	data cube, either
a)	by climbing up a concept hierarchy for a climension or
b).	by dimension meduction.
-	Fig (i) shows the result of a coll-up operation.
	performed on the central cube by climbing up
-	the concept hierarchy for location:
	The hierarchy was defined as the total order,
-	" steet < city < province _or_state < country"
	The coll-up operation shown aggregates the
	data by ascending the location hierarchy
->	from the level of city to the level of country. When real-up is performed by dimension reduction,
	one a more dimensions are removed from the
	given cube. The given cube.
	Eg: Consider sales data cube containing only the
	two dimensions location of time.
->	Roll-up may be performed by remains the time
	Roll-up may be performed by removing the time dimension, excelling in an aggregation of the
	total sales by location eather than by location
(avisa)	

	Page No 12
23/53	and by time.
(ii	Deill-down
and the same of th	de desdes production de la latera de latera de latera de la latera de latera delatera de latera de latera de latera delatera de latera de latera de latera delatera de latera de latera delatera delatera de latera de latera delatera delatera de latera delatera
-	Deill-down is the reverse of coll-up.
->	It navigates from less détailed data to
1:31-	more détailed data.
->	Duill-down can be realized by either,
aj.	Stepping down a concept hierarchy for a dimension
b).	introducing additional olimensions.
->	Figli) shows the result of a dill-down
	operation performed on the central cube by
23-1-2	stepping down a concept hierarchy from time
\rightarrow	defined as "day < month < quarter < year". Deill-down occurs by descending the time
	hierarchy from the level of quarter to the
unsti al	more détailed level of month,
->	Because a deill-down adds more détail to
	the given data, it can also be performed
	by adding new dimensions to a who
	Eg: a Suill-down on the central cube of fig (i) can occur by inteoducing an additional dimension, such as customer group.
	o, Jig () can occur by inteoducing an additional
,	aimenejon, such as customer group.
	Slice and Dice
	The transfer of the same of th
->	The slice oberation pulling 1 1-1:
op had	The slice operation performs a selection on one dimension of the given cube, resulting in
ad a	a sub whe.

	Proge No. 13 Dete 1
->	Fig(i) shows a slice operation where the sales
	the selected from the central whe for the
	dimension time using the witerion time = "g,"
-	The dice operation defines a subcube by performing a selection on two or more dimensions.
-	Fig (i) shows a dice operation on central cube
	based on the following selection criteria that
	involve the dimensions
anismerick	1 location = "Tounto" ou "Vancouvee") & (time = "9," a
_	"Q2") 4 (rtem = " home entertainment or computer)
	the shows the south of a dill-double
- iv	Pivot (xotate)
-	Dit o proposition to the state of the state
	Pivot is a visualization operation that estates the data ones in view in order to provide
	an alternative presentation of the data.
-	Fig(i) shows a pivot operation where the item
- 4 15	l hocation ances in a 2-D slive are
30000	eotated.
	tides and engineers to as with
- 30 ad	MODULE - 2 3 - no nout this
- 1	the con the by wheelering an in
- 3/	Explain data cube computation and curse of
Andrew	dimensionality.
TINE	- Multidimensional data Analysis is the Efficient
400 00	Computation of Aggregation across many sets
	In SQL terms, these aggregation are referred as
	"geoup-by". Where each group-by can be
	0 0
	Coopped by Top Cooppe



3) 6) Explain different methods of indexing OLAP data.

Ans: - To facilitate efficient data accessing, most data une bouse systems support inder Stendures and materialized views there are two methods of indening OLAP, they are; Bitmap Indexing ii) Join indexing. Bitmap Indening: - is a popular indening method, because it allows Quick Searching in data cubes, In the bitmap index jou a given atteibute, there is a distinct bit vector (BV), for each value 'v' in the atteibute domain. If a given atteibute's domain consists of n values, then 'n' bits are needed for each entry in the bitmap indea. -) If the atteibute has the value v for a given eon in the data table, then the bit expresenting that value is set to "1" in coveresponding eow of bitmap index. All other bits for that now are set to 'o'. Eg: - "All Electionics" data warehouse ; Suppose the dimension "item" at the top levels has four values such as "Home Entertainment, "Computer", "phone" & "Security". Bitmap indening is advantageous compared to hash' or Tree indices. Berause comparison, Join & aggrégation operations are reduced to bit Aeithmatic

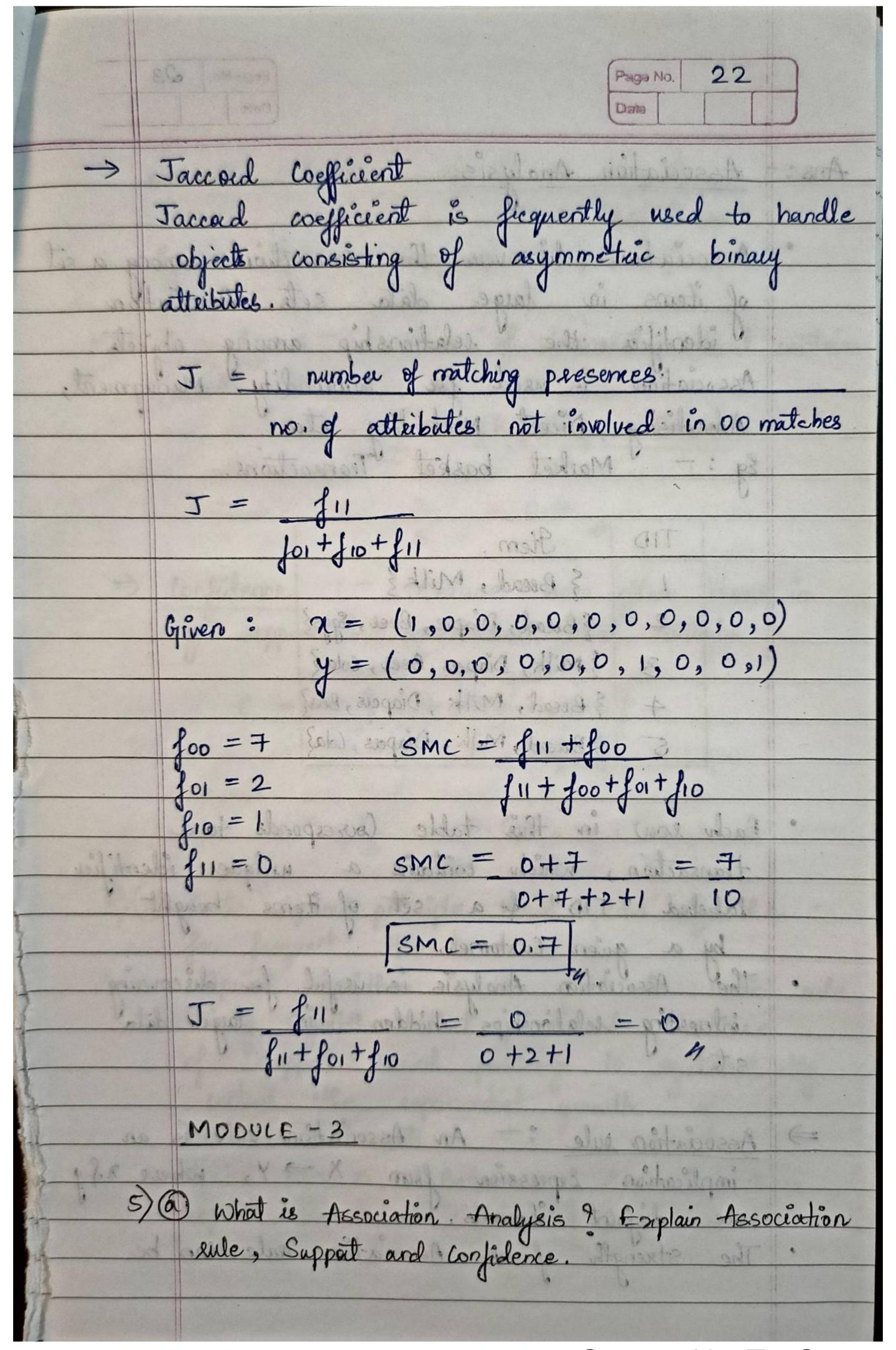
Peocessing time. Bitmap Indoxing provide or Significant reduct in space & input/output (I/O) since string of character can be represented by a single bit. Base table item bitmap table city bitmap RID item city RID H C P 3 RID V	tion					
Base table item bitmap table city bitmap						
RID HOM City RID H C P B RID V	table					
TO IVALLE	T					
RIH V RII O O O RII	0					
R2 C V R2 0 1 0 0 R2 1	0					
R3 P V R3 0 0 1 0 R3 1	0					
R4 5 V R4 0 0 0 1 R4 1	0					
R5 H T R5 1 0 0 0 R5 0	1					
R6 C T R6 0 1 0 0 R6 0	1					
R7 P T R7 0 0 1 0 R7 0	1					
R8 5 T R8 0 0 1 R8 0	1					
ii) Join Indering: Method gained popularity from its use in relational database query processing. Traditional Indering Maps the values in a given volume to a list of xoves having the values Join indering register the joinable lows of two relations from a relational database. If two relation R(RID, A) and S(B, SID) Join on the attribute 'A' & 'B', then the Join index record contains the pair (RID, SID), where RID & SID are record identifier from the R and S relations.						

	Prope No. 17
7	To of the sould be to the soul
and a	Join index seconds can identify joinable tuple without performing costly join operations.
->	Join indering maintain relationship between attribute
	values of a dimension (within dimension table)
	& the Couesponding now in the Fact table.
->	Consider &g: "All Electronics" Star Schema of a Join index
sidel- gove	relationship between the "Sales" fact table of the
TN	location '& item' Dimension table.
0	Sales 9
0 1	Res C V Proposition of the second
0	Ra O 1 O 1 O 1 S
	location T57
-	THE TO POLICE TO THE SECONDARY OF THE SE
•	Main Street T238 Sony-TV
1	
	T4598
. 9	17884
1	Asampa How Hold Hold Hold Co
	jig: Linkage between sales jact table & location &
	fig: linkage between sales fact table & location & item dimension table.
Salder S	able mineral
bocation	Sales-key item Sales-key location item Sales-key
1	CONTROL CONTRO
Main S	theet T57 Sony-Tv T57 Main Steet Sony-Tv T57
Main S	theet T238 Sony-TV T459
	treet T884
1 20 . 3	1991: 1: 1 hass so one of any one
	withdra a ten a sal

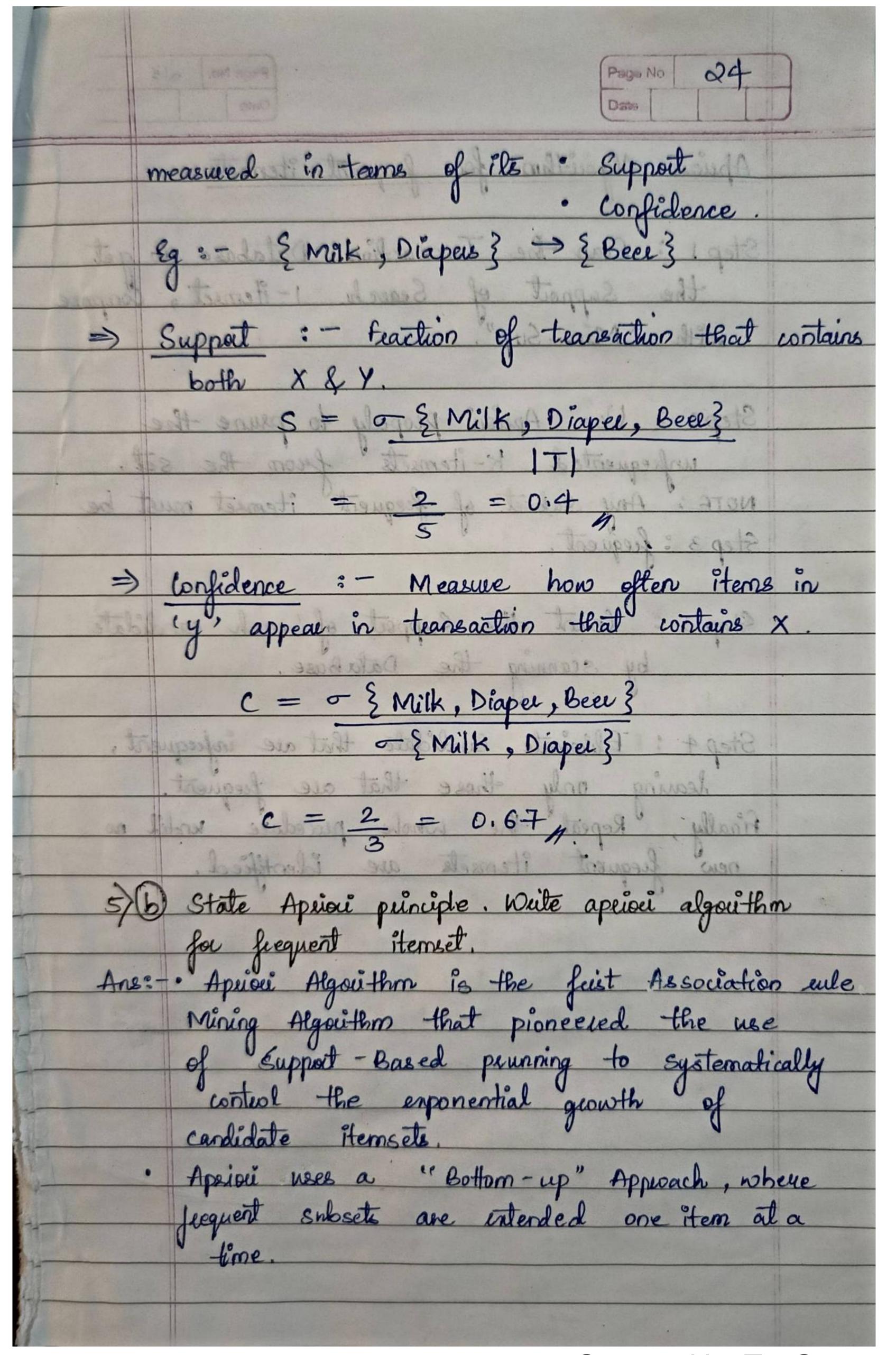
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chi	tee	10	Home	Marital	Annual. Income	Defaulted Borrower	
An	alysis	1 2	Yes No	Single	185 K	No.	Predictive Modelling
		3 4	No Yes	Single	70K 120K	No	X 0.
- Eisyland	Stadio City	5	No No	Disrocced	95K '	Yes .	
282	it with	7	Yes No	Divoced	220k 85k	No	
interest minis	astab a	9	No No	Married	75K	No Yes	
	Milk Milk	>6		diation lysis.	Anamoly Detection		
	1011.	*	dead	the coce	10 (10)	ning tasks	
	Paller	00	that the	sis :- desuibe data	stronge	y associa	ted.
	in the	ts.	form	patterns of implies	ation typ	eules or	presented feature

	Page No 20
	Page: No 20 Date
->	Because of the enponential size of its sough
	Because of the enponential size of its search space, the goal of association analysis
	is to enland the most interesting patterns
Onl-	in an efficient mannee.
Cin:	chiete A.C a a.
my.	of closely seeks to find groups
ase affan	chister Analysis: - It seeks to find groups of closely related observations, so that belongs to the same christer are more
	Similar to each other than observations
	that belong to other clusters.
→	Clustering has been used to group sets of
9 3	selated customers.
(vi locality)	Anamoly detection: - It is the task of
	identifying observations whose characteristics
le fiet	are significantly different from the rest of
	the data. Such observations are known as
3.31- au	anamolis de outlines.
	The goal of an anamoly detective algorithm is to discover the real anamolies of avoid
	falsely labeling normal objects as anomolies.
->	A good anamoly detector must have high
1	détection eate 4 lon false eate.
	Send = rumber of noticing attribute values
4/	(b) Define Similarity and dissimilarity between the objects.
	Find smc and Jaccord's coefficient of two binary vertous.
	X = (1,0,0,0,0,0,0,0,0)
	Y = (0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0

	Proge No. 21 Dete Dete
Ans:	Similarity
	Similarity between two objects is a numerical
	measure of the degree to which the two objects are alike.
- sque	Similarity is higher for pair of objects that
23000	Similarities are usually non-negative & are often between 0 and 1.
	delaids soft at probabilities
	Dissimilarity of his middle of the second of
10	Dissimilarity between two objects is a numerical measure of the degree to which the two
	Objects au alike
10 100	Dissimilarity are lower for more similar pair of objects.
um Al-1	Dissimilarity fall in the interval [0,1] or in the lange [0, \interval [0,1] or in the
Mich	all distances the sale arealist it a
30	Simple Matching Coefficient [SMC].
3	It counts both presence and absence equally.
	SMC = number of matching atteibute values
11 100	number of atteibutes
	post of the same of the same
	SMC = \$11+ \$00
	for+fro+fro
	110,0,0,0,0,0,0)=Y



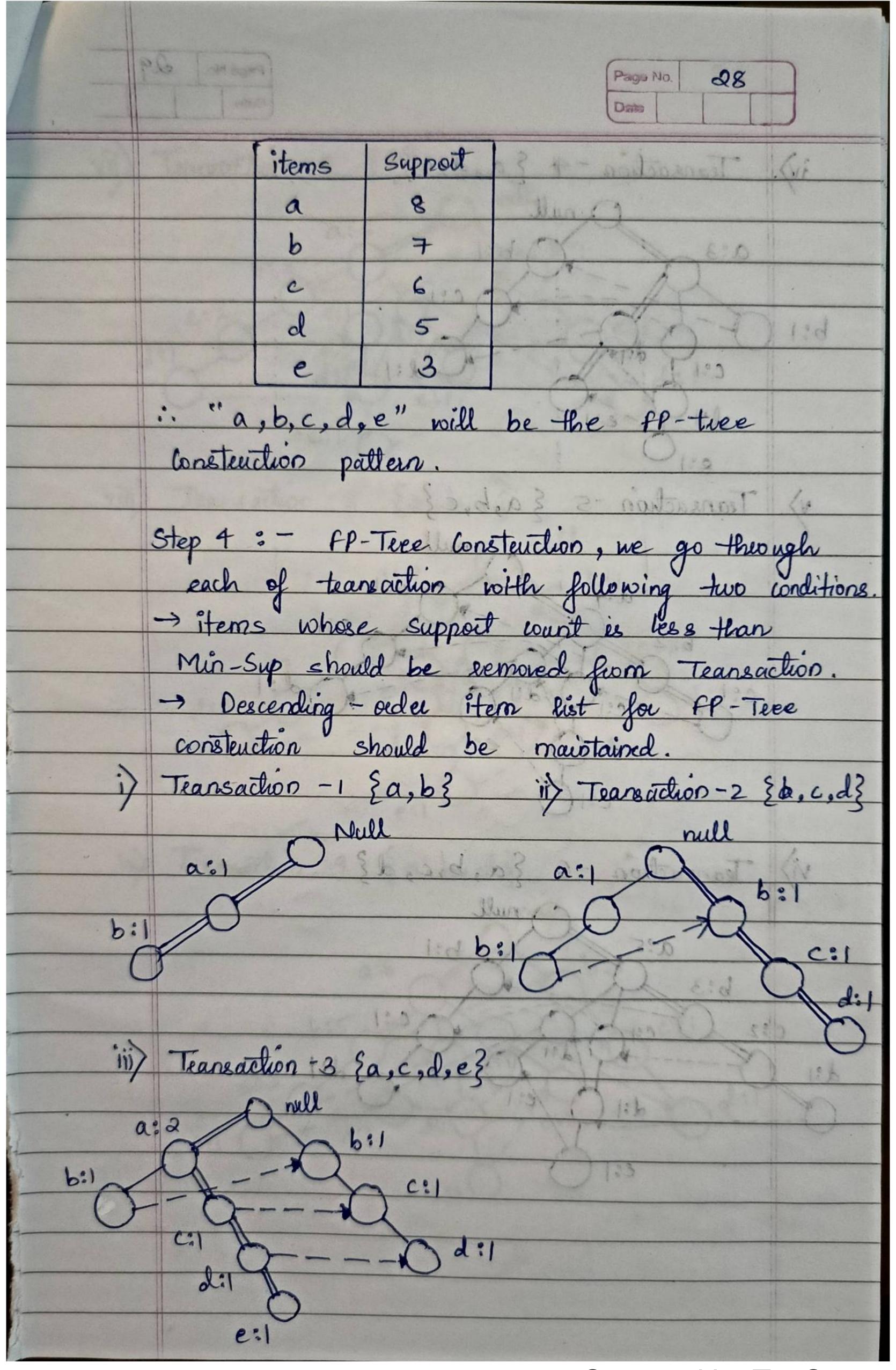
	Proc No. 23 Date
Ans:	Association Analysis
sthud	of hear alternat is to signed heart
•	Association discovers the connection among a set of items in large data sets. Association
	identifies the relationship among objects.
· · ·	Association is used for commodity Management,
236.59.00	Advertising, Direct Marketing etc. Eg: - Market basket Transactions.
	Eg : Markel baskel lansactions.
	TID Stem.
	1 3 Bread, Milk 3
	2 ? Bread, Diaper, Beer, Eggs?
	3 { Milk, Diapers, Beer, Cola}
	4 & Bread, Milk, Diapers, Béer?
	5 & Buead, Milk, Diapers, Cola}
	101 + 101 + 100 + 101 + 100
•	Each eon in this table Corresponds to a
	teansaction, which contains a unique identifier
	labeled "TID" & a set of items bought
	by a guien Customer.
	The Association Analysis is useful for discovering interesting relationships hidden in large data
	sets, delationships hidden in large data
	01, 10, 11,
=	Association rule: An Association rule is an
	implication enpression form X -> Y, where nfy
ALTERNATION OF	and susjoint Hemselb,
•	The strength of an Association rule can be

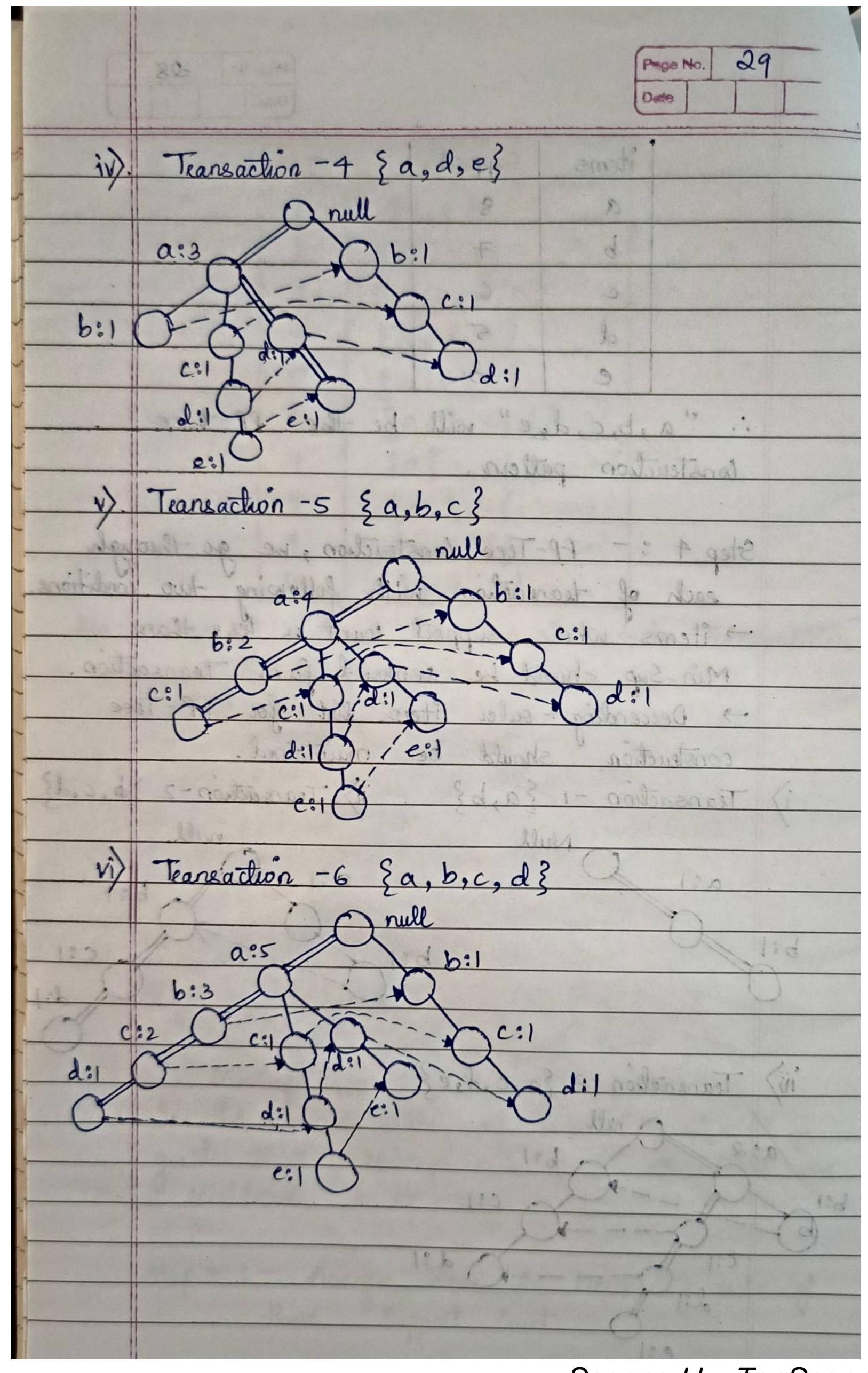


	Pmgo No.] 25
	Dute
	Aprion Algorithm
12.	K=1
2).	Find all frequent 1-items ets fx = {ili & I ^
1	$\sigma(\{i\}) \geq Minsup\}$
3).	k = k+1
4/.	
6	Fix = Apriori - gen (Fx-i) //generate Cardidate Hemset. For each teansaction teT do
	Generate Subset (CK, t)
	End foe
9).	Fx = { c c e cx > minsup } // Extent K- heavent
	Fx = {c ceck > Minsup} // Extend K-fleggent itemset.
10).	until $f_K = 0$.
W).	ection FK.
60	:- TID Hems
7	3 Buad, Milk?
	2 {Bread, Diaper, Beer, Eggs}
	3 {Milk, Diaper, Beer, Cola} MinSup Court = 3.
	4 & Bread, Milk, Diapee, Beaf
	5 3 Bread, Milk, Diaper, Cola ?
-	Firstly, Generate 1-itemset by counting the components.
1	Bread 4
	Diapee 4
	Milk 4
	Beer 3 (1-itemset)
	1/299///
	Cooperad by Top Coopera

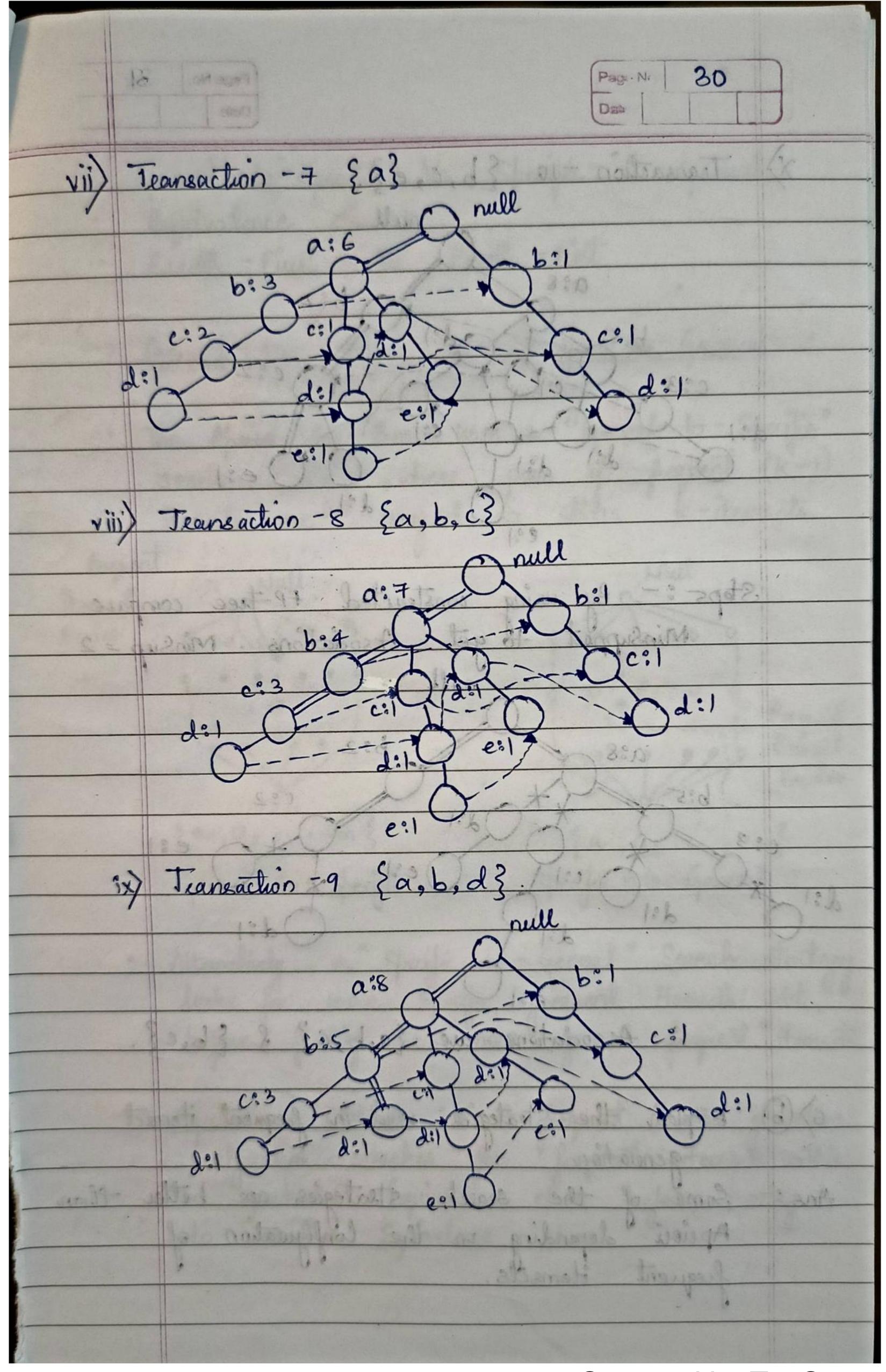
	Page No. Q6	
	Eleminate the candidate that are infrequent by comparing with 'MinSupCount = 3' such 'Egg & Cola'	
	Next, Generate 2-items et by forming sub of two components from 1-items et.	set
	Ebread, Milk 3 Ebread, Diaper 3	
	[2-itemset). Milk, Diaper? 3	
The second of	/{ Milk/, /Beel? //2// { Diapu, Beel} 3	A
	Get the count of Subset's by scanning Bows Eliminate the candidates that have less	e-DB,
->	Generale 3-itemset by following subset of the components from 2-itemset.	vee
	Ebread, Milk, Diaper 3	
	Ebread, Diaper, Beer 2 (3-itement)
3.03	2 Milk, Diaper, Beee 3 2	•
-	E Beead, Milk, Beer 3 1	
	All 3-itemset ave infrequent, Because count à	· .
23	less than Min-Sup Count = 3.	
6)6	6) Constant an FP tree for the following datase	₽ .
10	steps :- Avage Hens in Descending gude	
	ties trappet taint	

	Proce No. 27 Date
· Jan	TID Hems
an a	Sabbille pier pier
	2 \ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\
	3: {a,c,d,e}
	4 zagd, ez
	5 {a,b,c}
	6 {a, b, 0, d}
	7 {a} ogoid , bosses?
	5 2a, b, c3 bosses
	9 E{a,b,d} : xint?
	10 {b,c,e}.
Ans:	- Step 1: - Count all item in whole Transactions
Perse - 199-	to get support count.
- 21	Eliminate the foods after that I have to
	items Support
	Generale 3 stempt 6 shores
	c 6 c c c
	d 30500 11 M. bos 33
(4-	es (3 13 13 13 13 13 13 13 13 13 13 13 13 13
	Step 2: - Apply the Threeshold (MinSup) & comove
	item which posses Support court less than
- 14	MinSup.
-	In the above table all items Support count is
+	greater than "MinSup = 2", So all items are considered in FP-tree.
-	considered in FP-tree.
1	
-	Step 3: - Amange items in Descending order of
	there Support Court.
and the same of th	

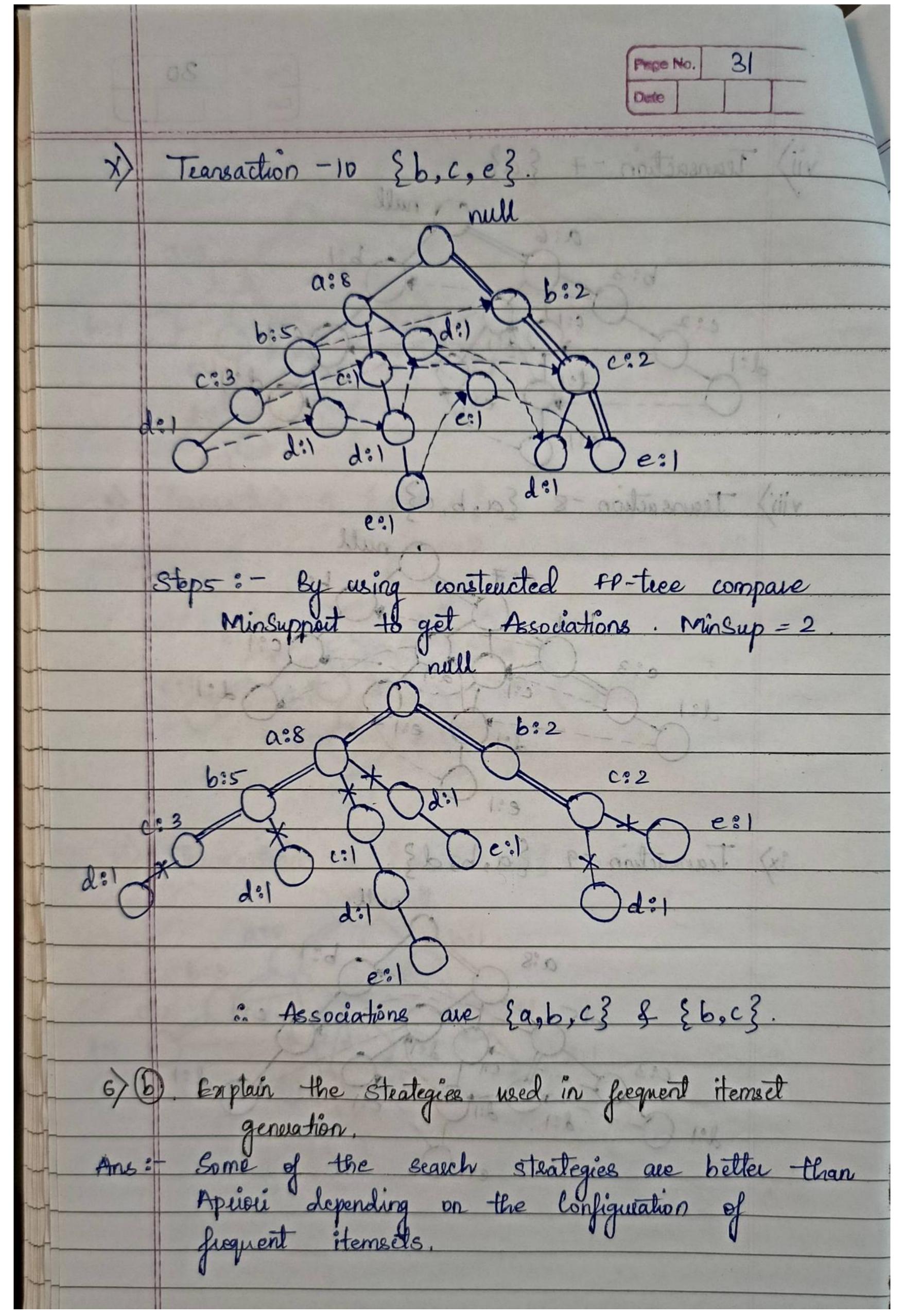


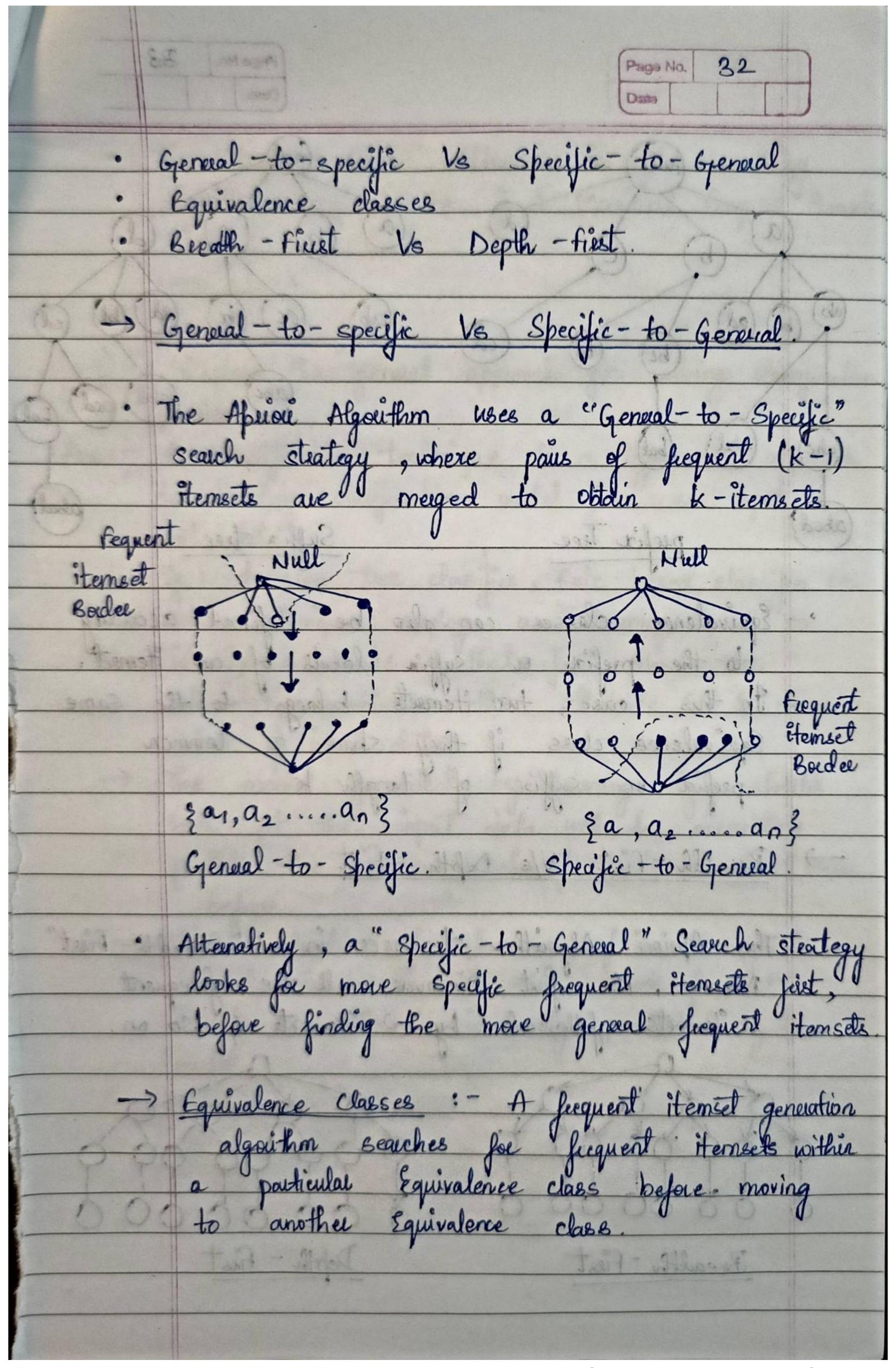


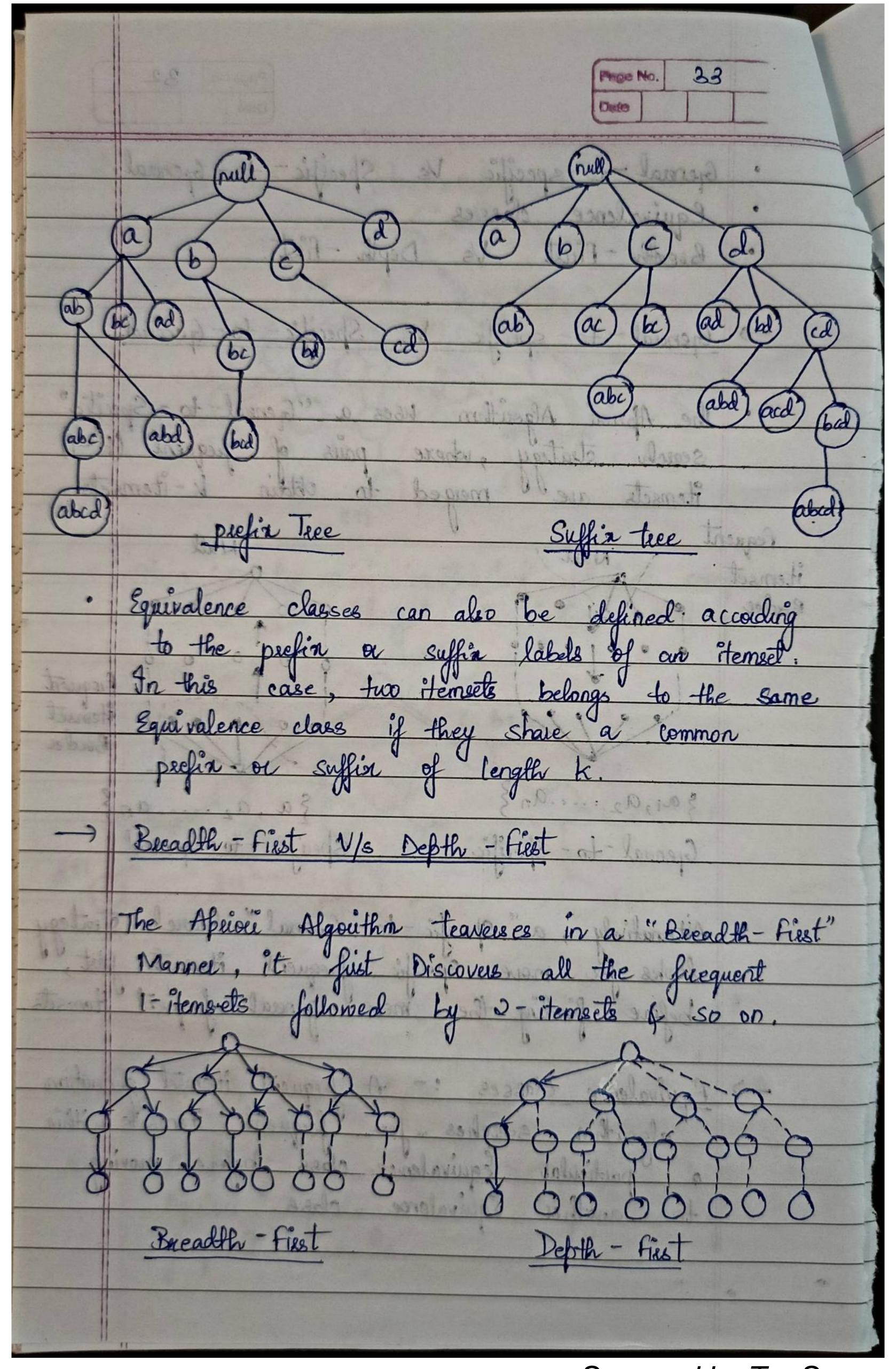
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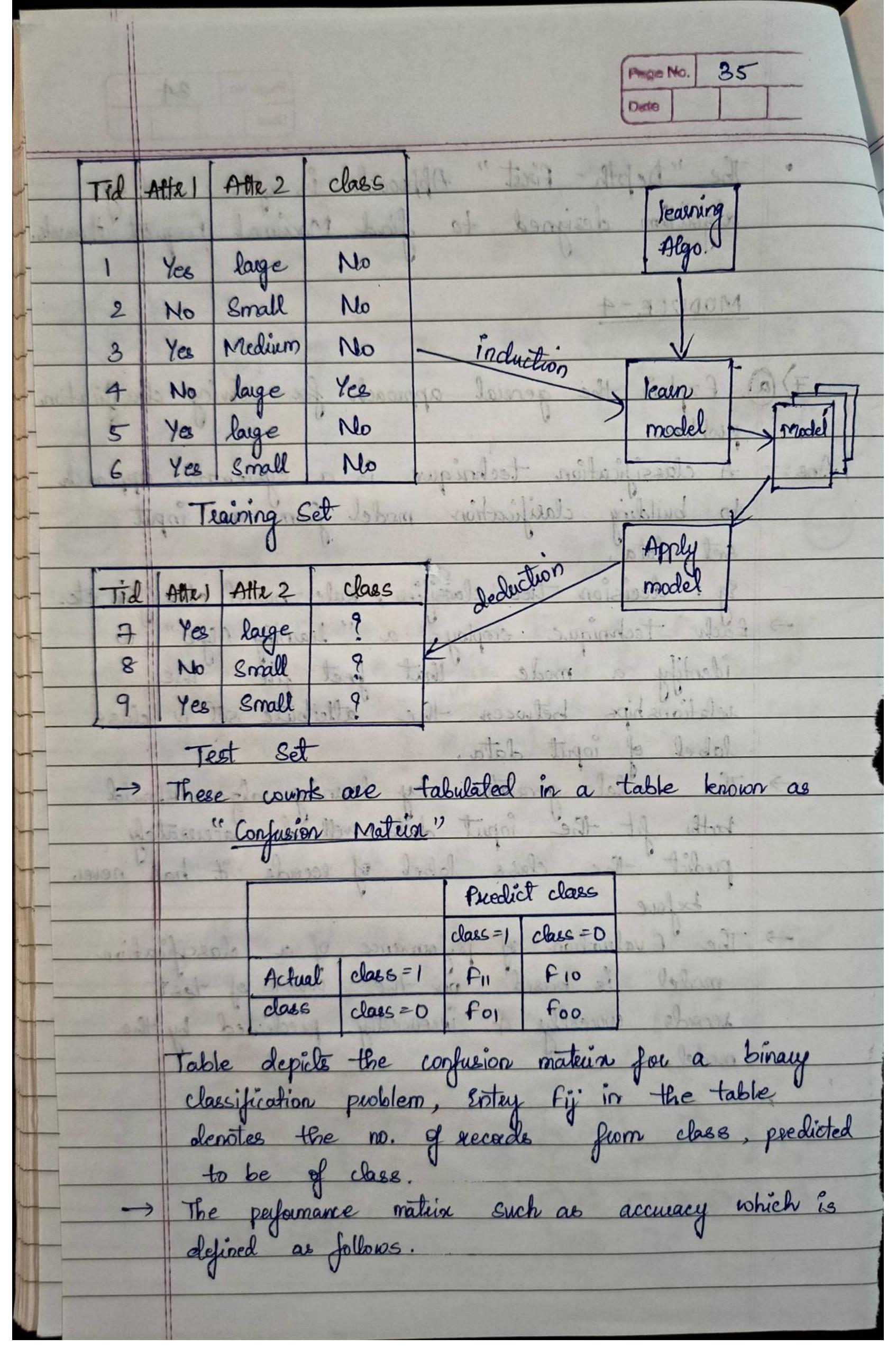
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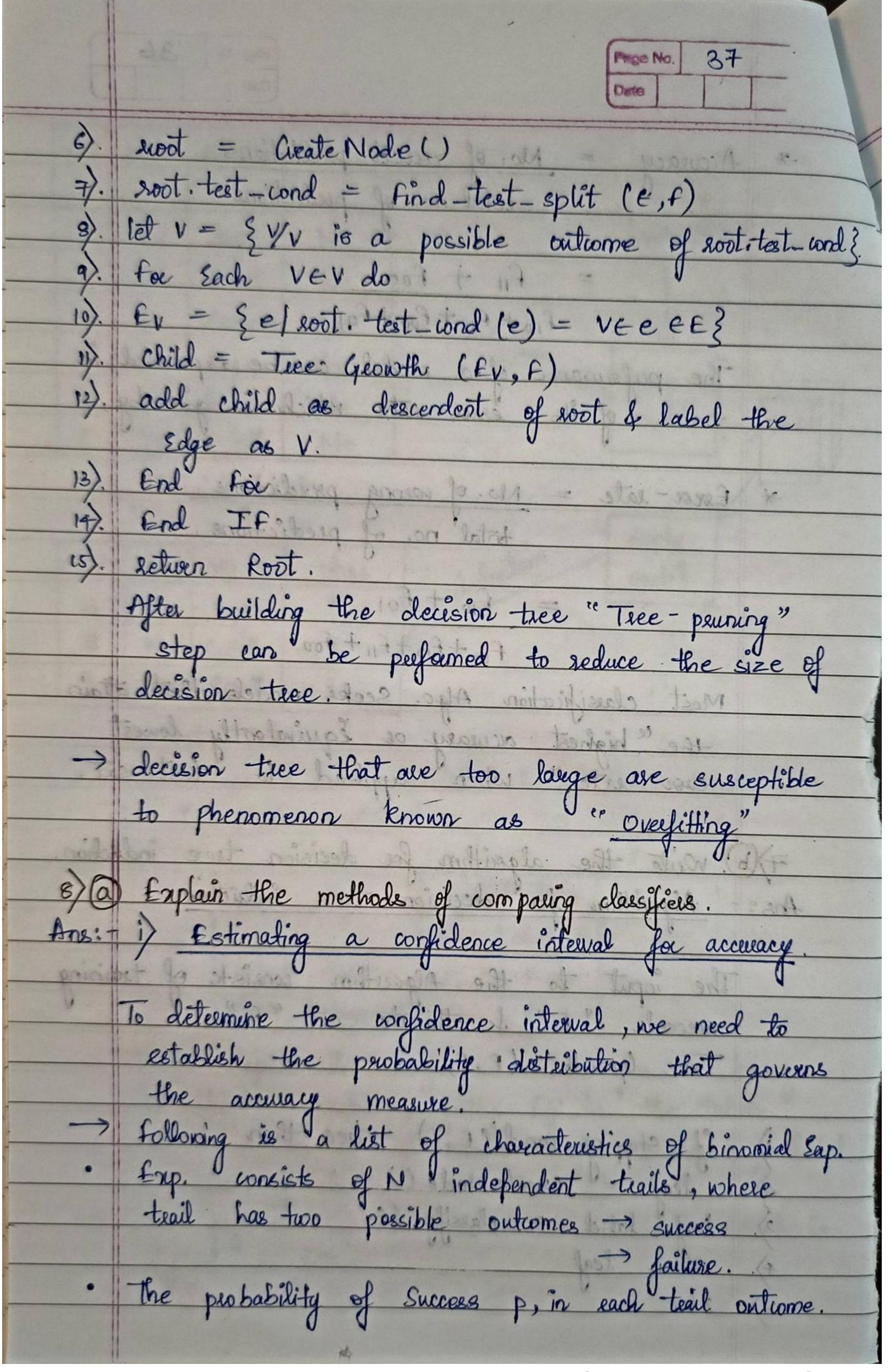


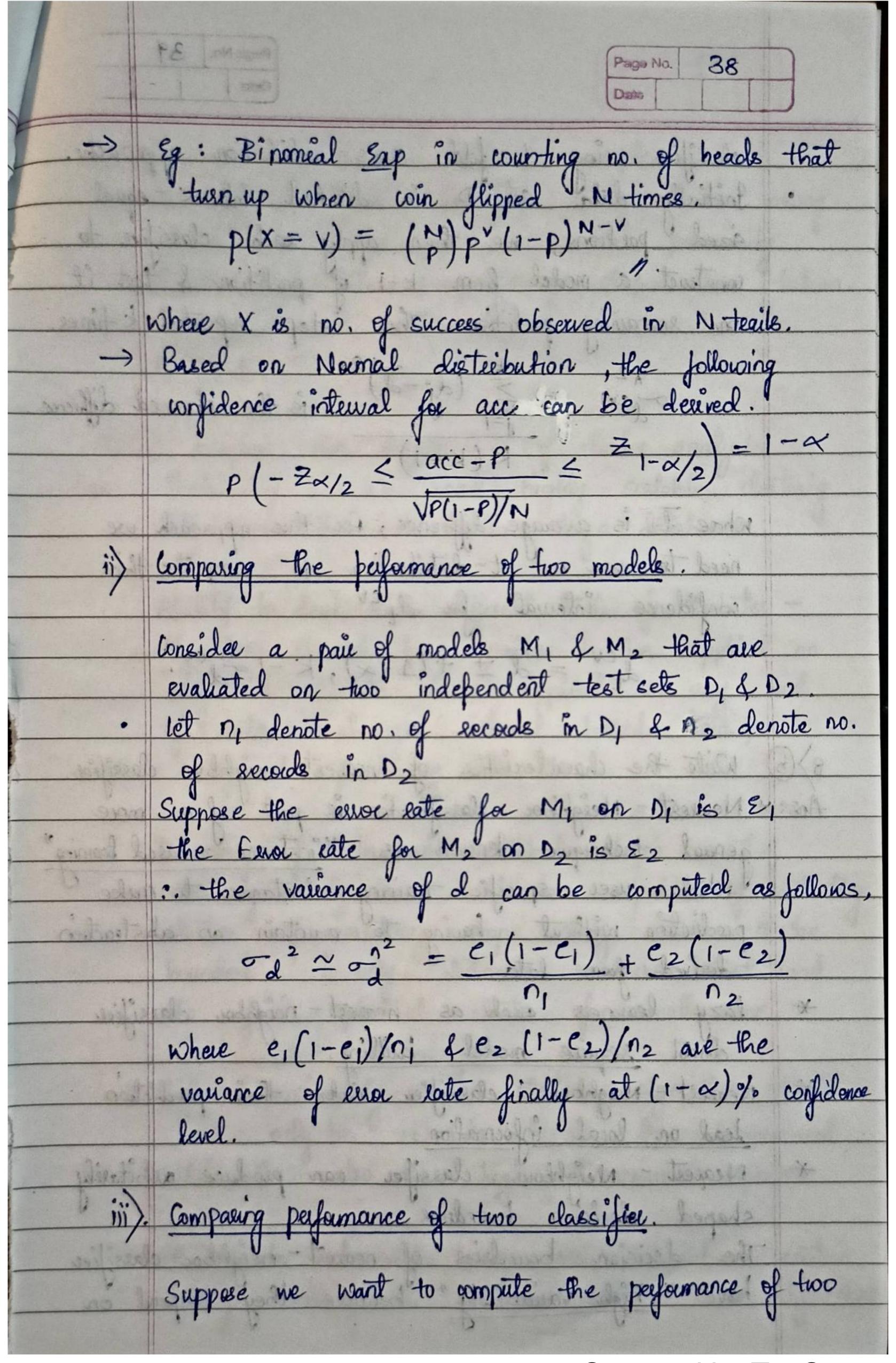


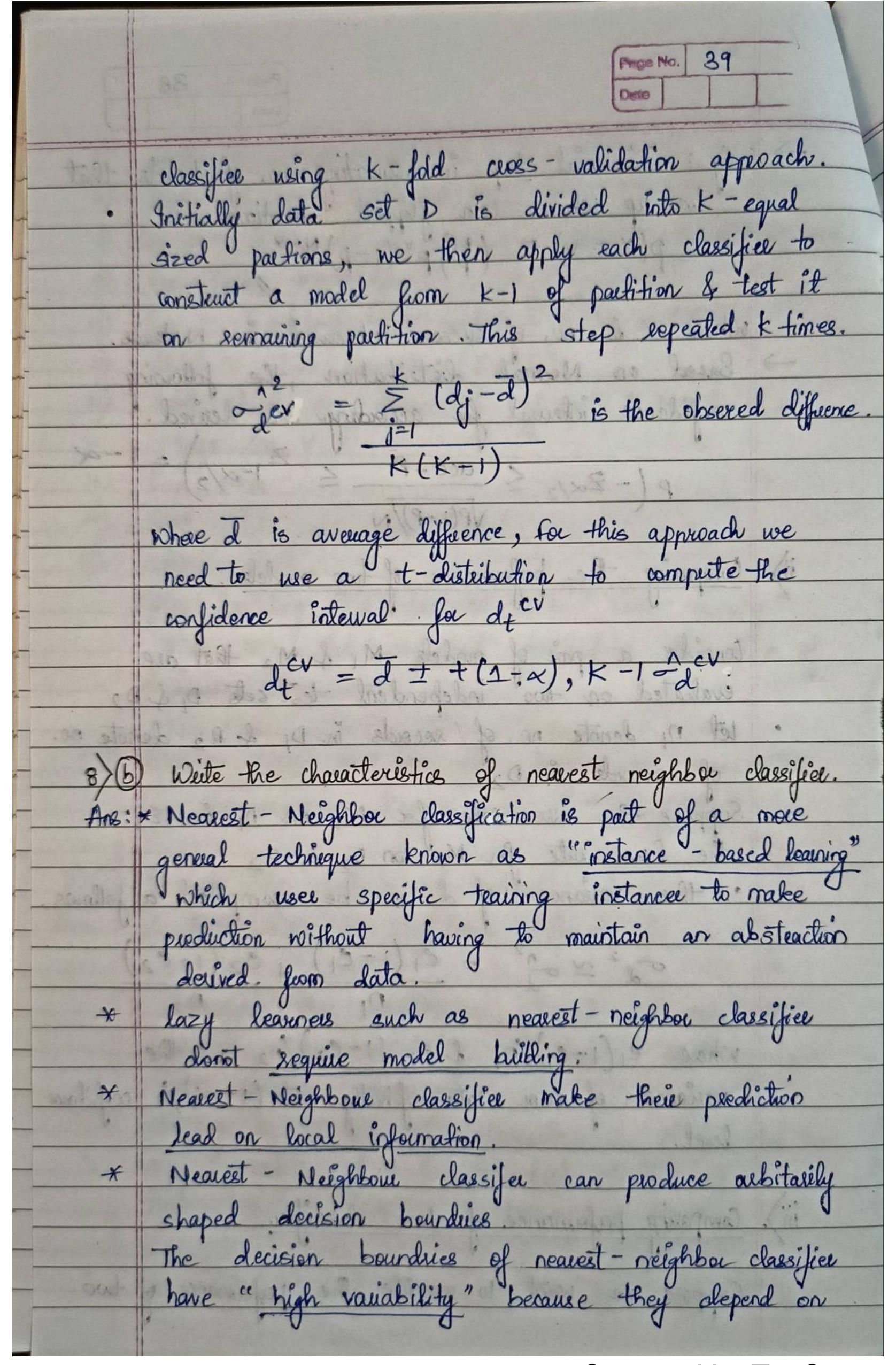
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36 Mo, of coceet predictions total no. of predictions = Fil + Food Vav book soll 6 F10+ F11+ F00+ F01 The performance of a model can be expressed in terms of its error-eate which is given by, * Essec-eate = No. of weong predictions total no. of predictions. = first for at positive inter fortfirt for Most classification Algo. Seeks models that attain the "highest accuracy or Equivalently lowest evor sate" when applied to test set. 7) (b). Write the algorithm for decision tree induction. Ans: - Algoeithm for decision tree induction. The input to the Algorithm consists of training execorde "E" of attendite set "F". Tuce Geowth (E,F) the occupacy measure if stopping-cond (E, F) = tree then leaf = create Node () leaf. label = classify (E ecturon legs

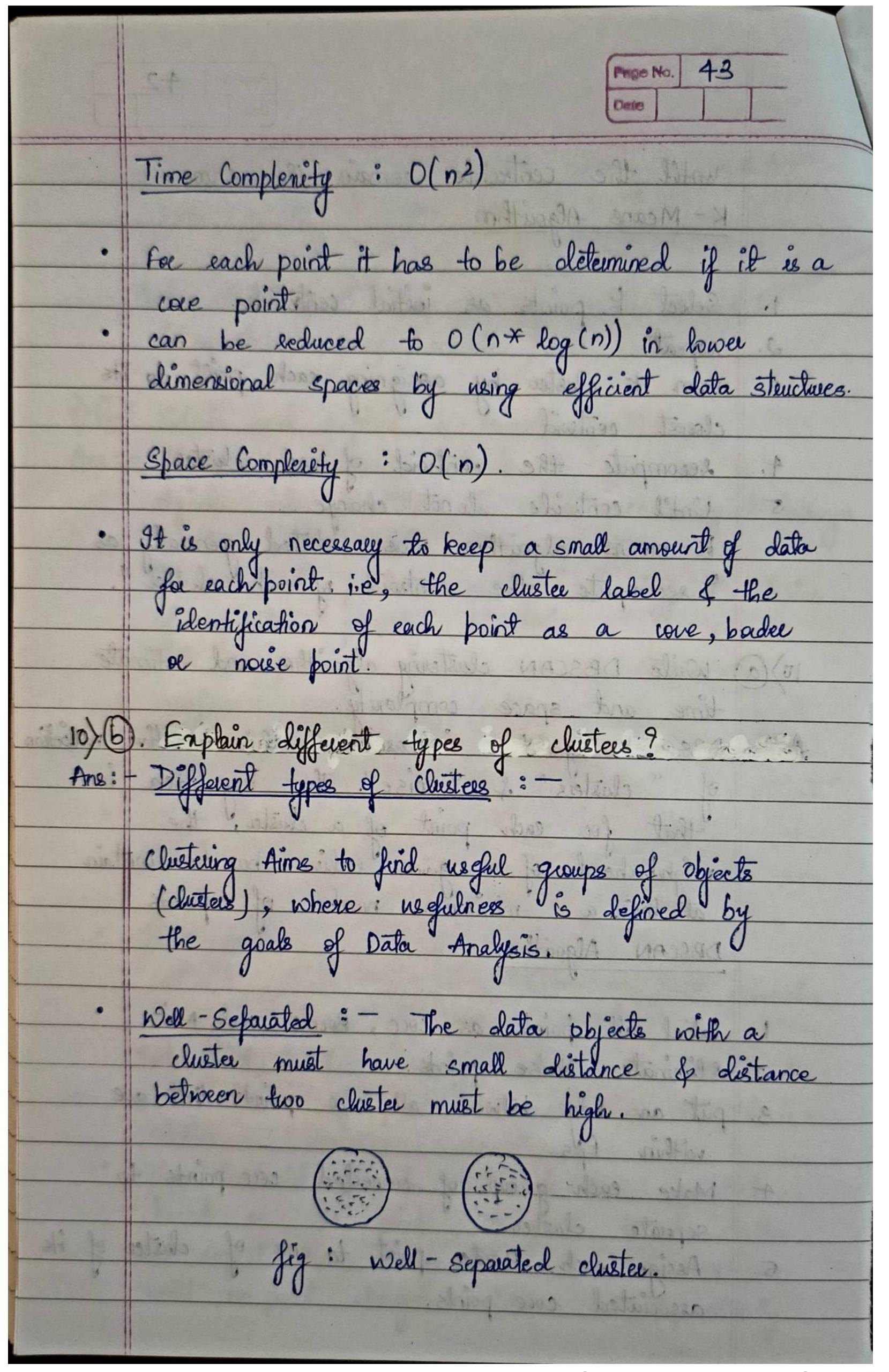


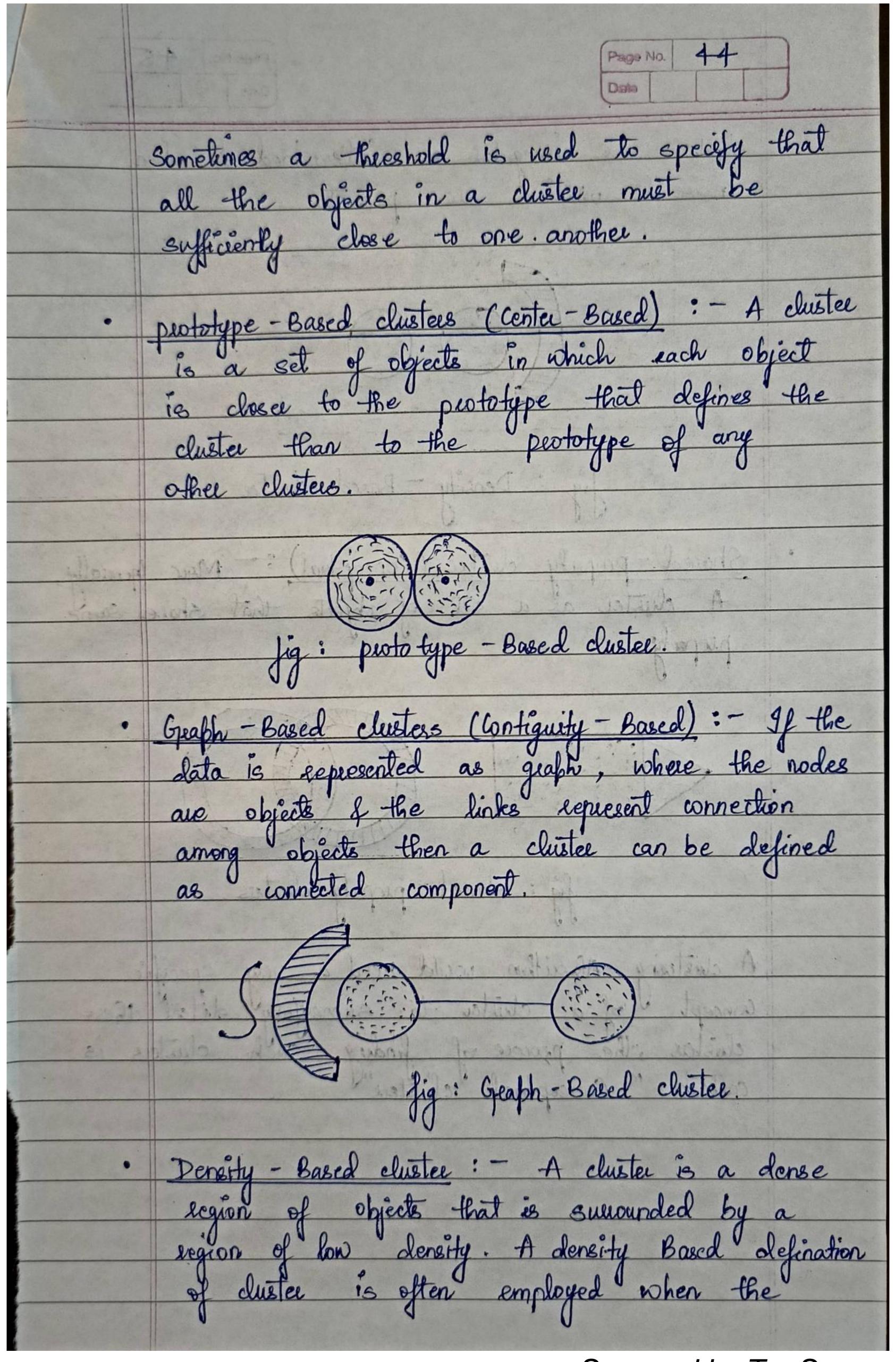


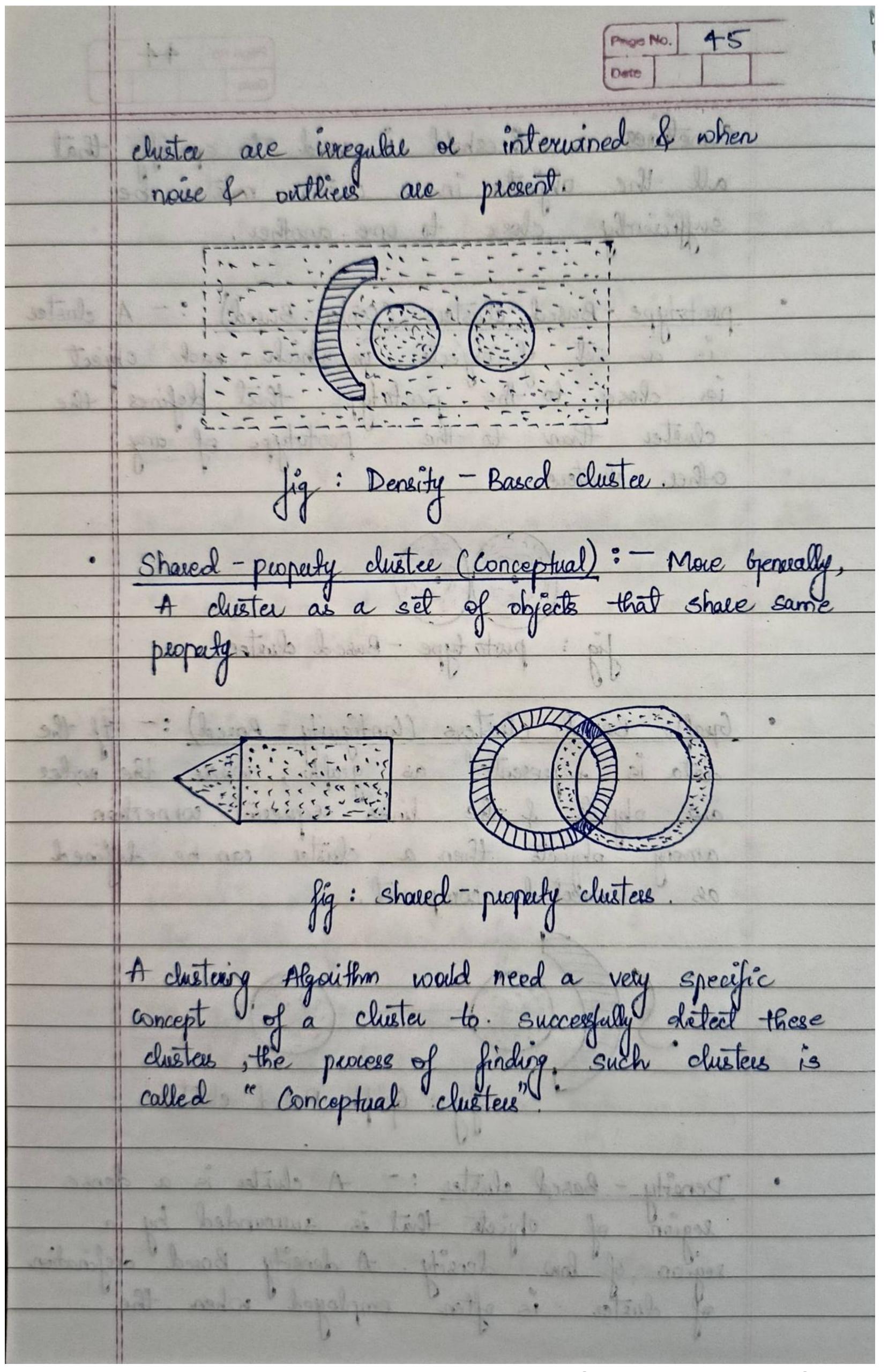


are sensitive to such data and may lead to poor qualify christers. Interpretability - The clustering exults should be interprétable, comprehensible, and usable. 9 (b). State and explain K-means algorithm. Ans: - K-means clustering intends to pactition 'n' objects into k-duster in which each object belongs to the cluster with the nearest mean. This Method produces enactly "k" different cluster
of greatest possible distinction. It is a prototype-based clustering technique creates a one-level partitioning of data objects. There are no. of such techniques but two of most perminent are "K-Means" & "K-Mediad" The objective of k-mean clustering is to minimize total "intra-cluster variance" or "squared error function ". In k-mean, we first choose "k" initial centroids, where "k" is a usei-specified paramèter, namely the no. of clusters desired. Each point is then assigned to closest conteoid of each collection of points assigned to a centerial is a chieter. The Centerid of each christer is then updated based on the points assigned to the cluster we repeat the assignment of update steps no point changes chieters or Equivalently

	Page No 42 Date Date
	until the centroide remain the same. K-Means Algorithm
	K- Means Higharm
	Select k points as initial centroids.
3.	comate (1) and the 10 of headless of one
3.	form k cluster by assigning each point to its
	closest centerid.
4.	ecompute the controid of each chuster.
5.	Datil centroide do not change.
and the same	1- moone algorithm was stated generally as
	execompete the controid of each cluster".
35%	the second standard to
10)(0	Wite DBSCAN clustering algorithm and estimate
	time and space complonity. - DBSCAN algorithm is based on this intuitive notion.
Hine:	of "clusters" of "noise". The key idea is
	that for each point of a cluster, the
	neighborhood of a given eading has to contain
	at least a minimum number of points.
	DBSCAN Algorithm
	label all points as core, Border and Moise
	Eliminate Noise points.
3	put an Edge between all core points that are
4	- Make each group of connected core points to
	separate cluster.
5	Assign each Border point to one of christer of its
	associated core points.







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